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Foundations of Query Languages Summer semester 2010 April 28, 2010

2. Exercise Set: FOL and its Relationship to Databases

Exercise 1 (Modelling in FOL)

Consider the following database scheme $\{D, P, S, I\}$, where D, P, S are unary relational symbols and I is a 2-ary relational symbol. Let the database $(D^{\mathcal{A}}, P^{\mathcal{A}}, D^{\mathcal{A}}, I^{\mathcal{A}})$ be given. We interpret $D^{\mathcal{A}}$ as the set of exactly all points and straight lines. $P^{\mathcal{A}}(a)$ means that "a is a point". $S^{\mathcal{A}}(a)$ means that "a is a straight line". $I^{\mathcal{A}}(a,b)$ means that "a is a point, b is a straight line and a lies on b". Write down in first-order logic:

- a) For two points there is exactly one straight line on which they lie.
- b) On every straight line there lie at least two points.
- c) The intersection of two straight lines is exactly one point.

Exercise 2 (Query Reformulation)

Let

$$\varphi := \exists z (R(x,z) \land \exists w S(w,x,y)) \land \neg x = y$$

be given. Rewrite the query $\{\langle x, y \rangle | \varphi\}$ to an equivalent formula in relational algebra. Do you habe all information that you need in order to do so?

Hint: In order to do this exercise you may make further assumptions on the schema.

Exercise 3 (Query Satisfiability)

Is there a finite database such that the relational calculus query associated with φ yields a non-empty result? If so, give an example for such a structure (together with an underlying domain). If not, then prove your claim.

- $\varphi := \forall y (P(y,c) \Leftrightarrow P(y,y)) \land \neg P(x,x)$
- $\varphi := \forall y (P(y,c) \Leftrightarrow P(y,y)) \land \neg P(x,x) \land P(x,c)$

Assume that c is a constant.

Exercise 4 (Undecidability)

Let φ, ψ be first-order sentences as input. Conclude from Trakhtenbrot's Theorem that the following problems are undecidable:

- a) equivalence of φ and ψ ,
- b) containment of φ in ψ , and
- c) domain independence for input φ .

Due by: May 5, 2010 before the tutorial starts.

Literature: S. Abiteboul, R. Hull, V. Vianu: *Foundations of Databases*, Addison-Wesley, 1995. ISBN 0-201-53771-0. Download available at http://www.inf.unibz.it/~nutt/FDBs0809.